

On the Efficient Evaluation of Array Joins

Big Data in the Geosciences Workshop
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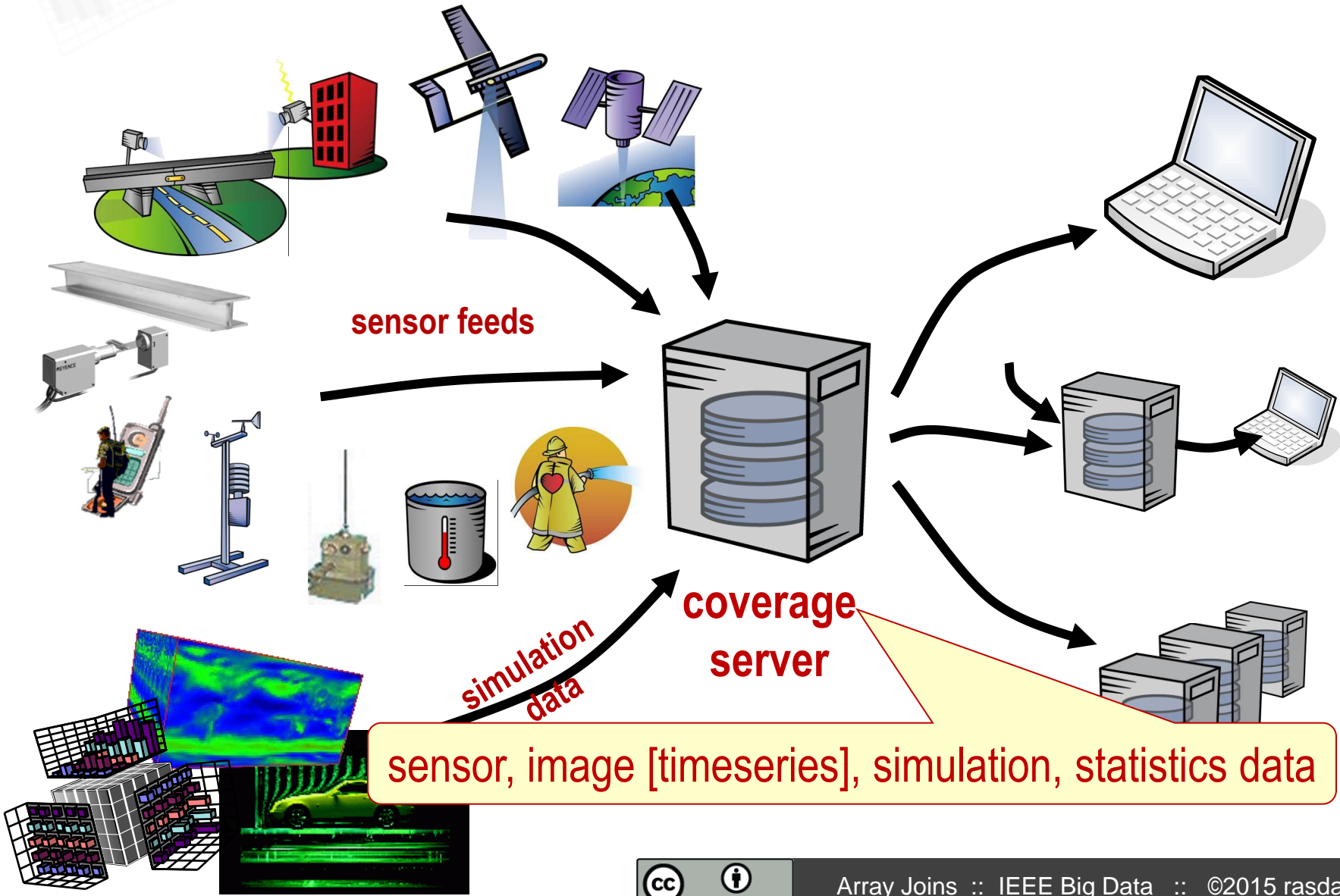
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[gamingfeeds.com]





Data Homogenization With OGC Standards




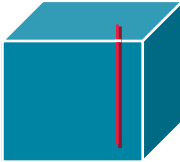
Web Coverage Service (WCS)

- OGC **Coverages** unifying regular & irregular grids, point clouds, meshes
 - OGC Coverage Implementation Schema
- **WCS Core**: access to spatio-temporal coverages & subsets
 - subset = **trim** | **slice**





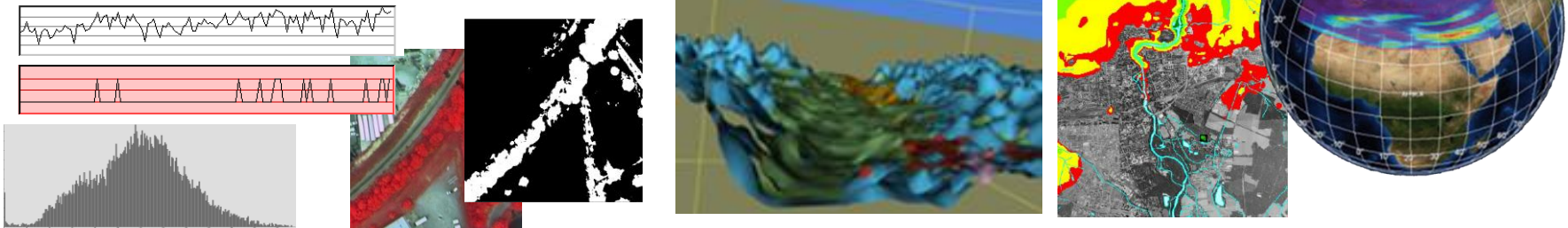



- **WCS Extensions**: optional functionality facets
 - Scaling, CRS transformation, ...

Large, growing implementation basis:
 rasdaman, GDAL, QGIS, OpenLayers, OPeNDAP, MapServer, GeoServer, NASA WorldWind, EOx-Server; Pyxis, ERDAS, ArcGIS, ...

OGC Web Coverage Processing Service (WCPS)

- = high-level spatio-temporal geo analytics language std



[JacobsU, FhG; NASA; data courtesy BGS, ESA]

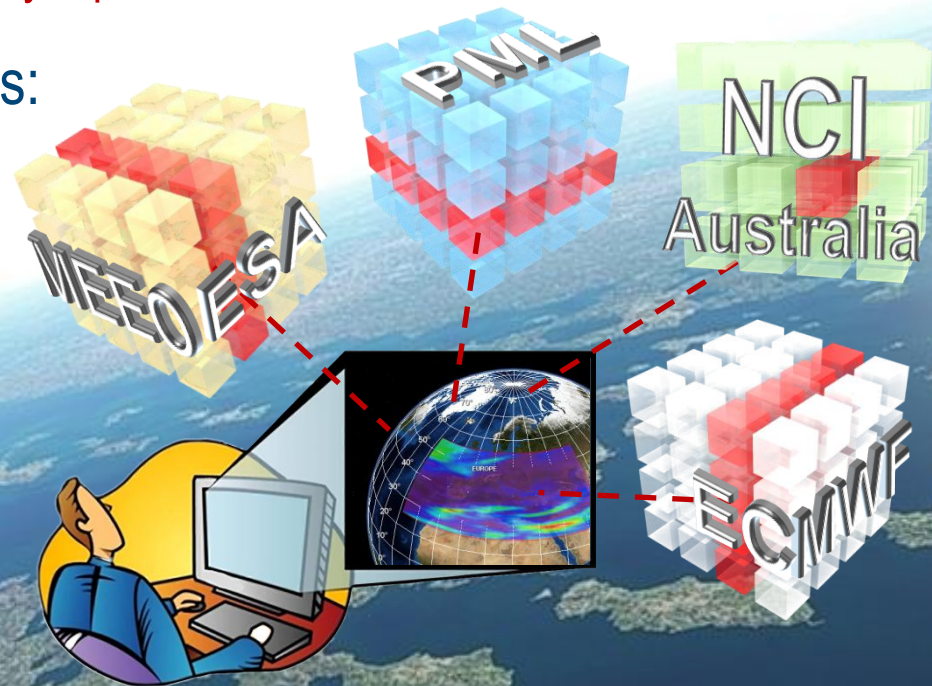
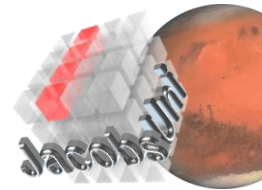
- "From MODIS scenes M1, M2, M3: difference between red & nir, as TIFF"
 - ...but only those where nir exceeds 127 somewhere, within LandMask

```
for $c in ( M1, M2, M3 ),
    $lm in ( LandMask )
where
    some( $c.nir > 127 and $lm )
return
    encode( $c.red - $c.nir, "image/tiff" )
```



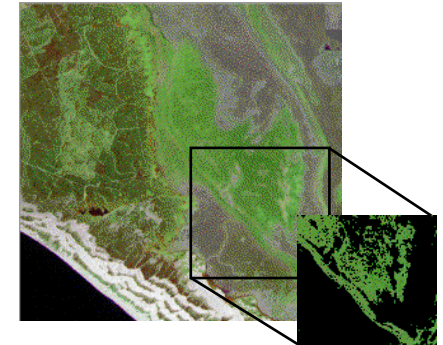
EarthServer: Datacubes At Your Fingertips

- **Agile Analytics** on Earth & Planetary **datacubes**
 - rasdaman + NASA WorldWind
 - Rigorously standards: OGC WMS + WCS + WCPS
 - 100s of TB online now, next: **1+ Petabyte per cube**
- Intercontinental initiative, 3+3 years:
EU + US + AUS



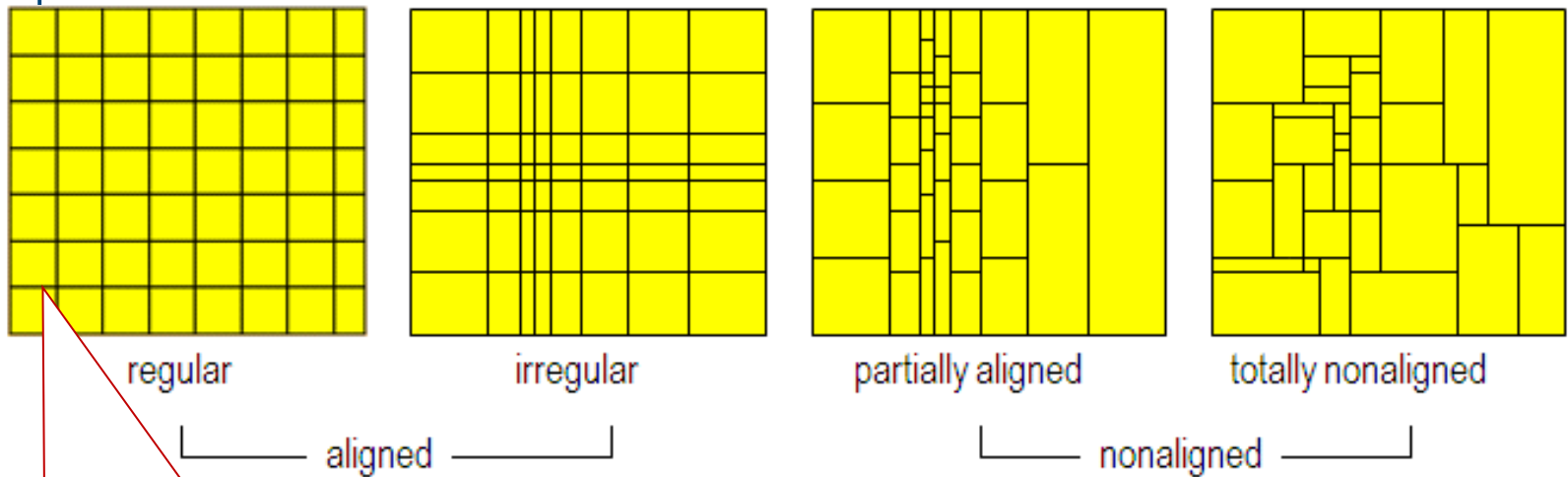
Agile Array Analytics: rasdaman

- „raster data manager“: **SQL + n-D arrays**
- Scalable parallel “tile streaming” architecture
 - Joins!
- Supports R, QGIS, OpenLayers, MapServer, GDAL, EOxServer, Pyxis, ERDAS, ArcGIS, ...
- Blueprint for OGC WCPS, ISO Array SQL stds



Array Partitioning

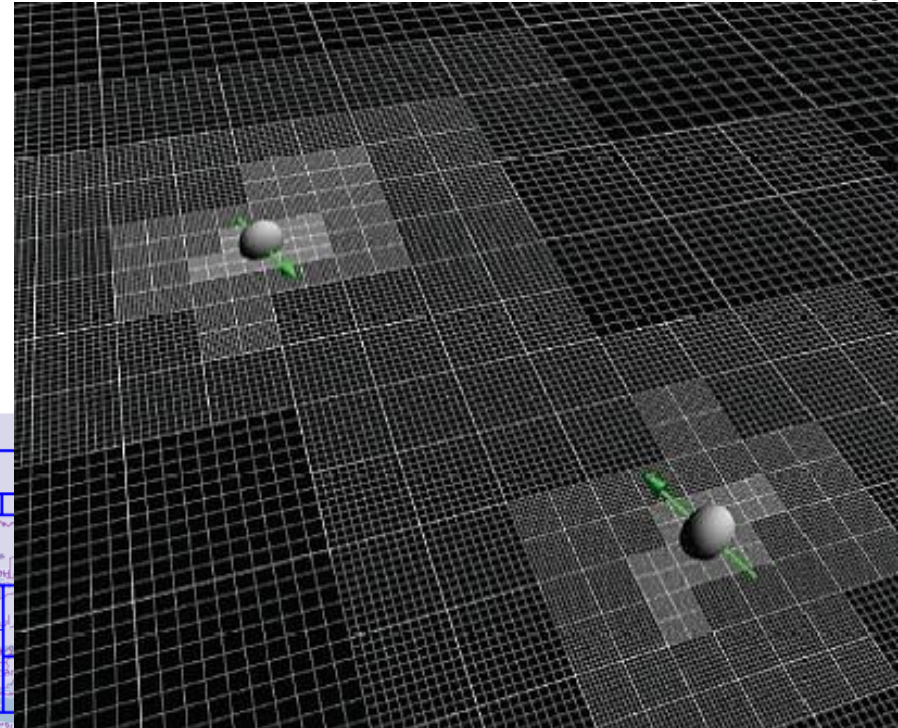
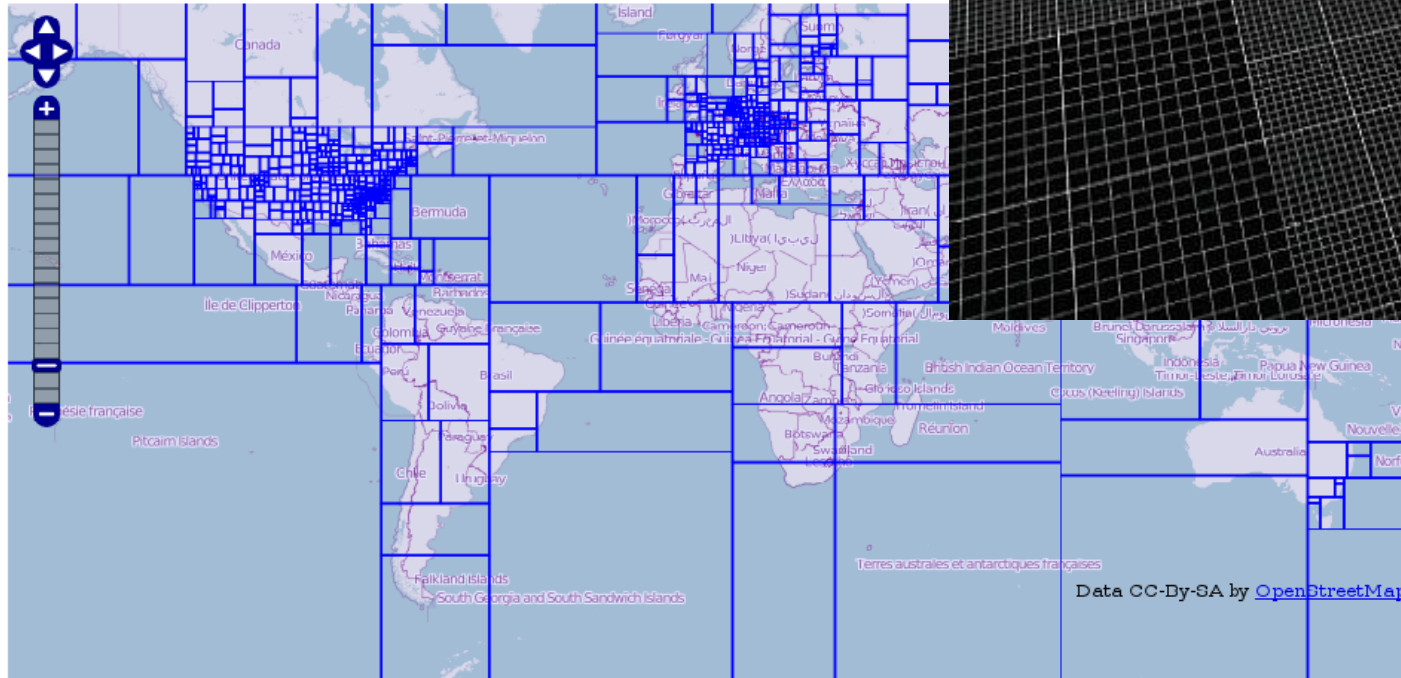
- Goal: **faster loading** by adapting storage units to access patterns
- Approach: split n-D array into n-D partitions („tiles“)
- Tiling classification based on degree of alignment [ICDE 1999]
 - all implemented in rasdaman



chunking [Sarawagi,
Stonebraker, DeWitt, ...]

Why Irregular Partitioning?

[Centrella et al: scidacreviews.org]



[OpenStreetMap]

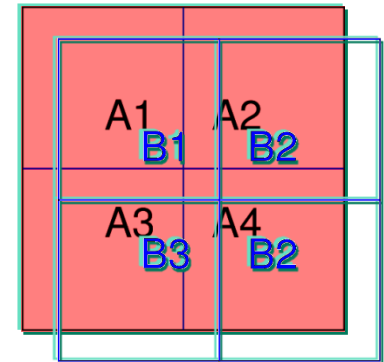


Array Join: What Happens?

- „MODIS red band with cloud mask applied“:

```
$Modis.red * $CloudMask
```

- Case 1: same tile shapes, same position
 - Easy
- Case 2: same tile shapes, different position
 - Overlaps, not so easy
- Case 3: different tile shapes
 - Worse
- Case 4: different dimensions
 - Gimme a break!



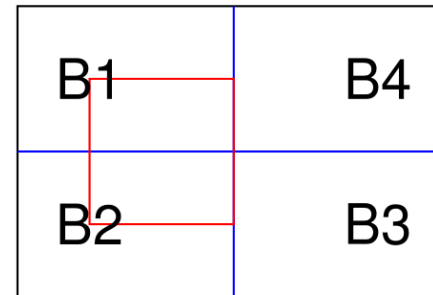
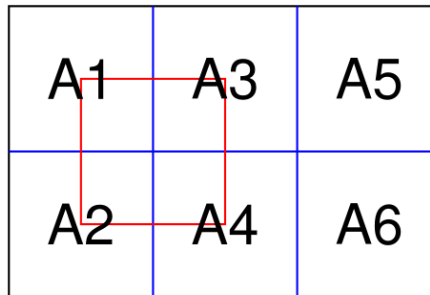
A1	A3	A5
A2	A4	A6

B1	B4
B2	B3

Array Join: Problem Statement

- Goal: minimize tile reads when evaluating „A op B“ in face of some arbitrary, independent partitioning of A and B

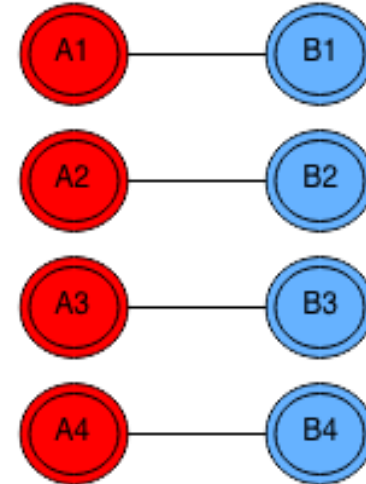
`$Modis.red * $CloudMask`



Bipartite Traversal Graphs

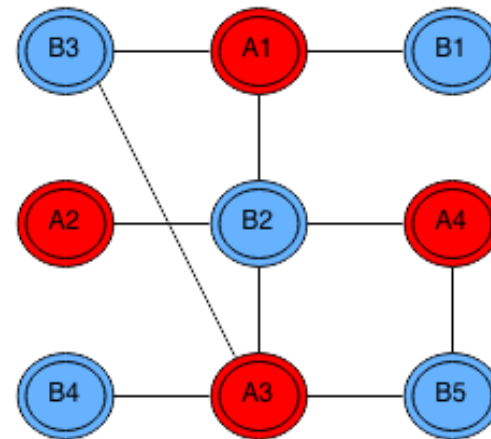
A1	A2
A3	A4

B1	B2
B3	B2



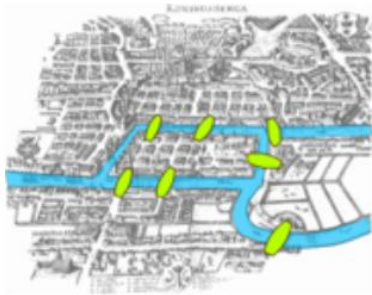
A1	A2
A3	A4

B1	B2
B3	
B4	B5

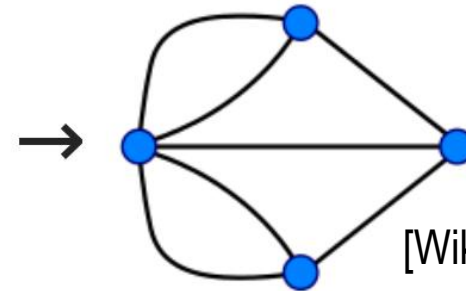
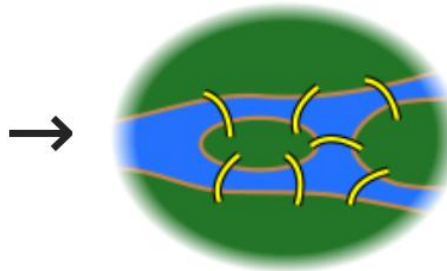


Finding Complete Paths

- Leonhard Euler: „complete, minimal path for **even-degree nodes**“

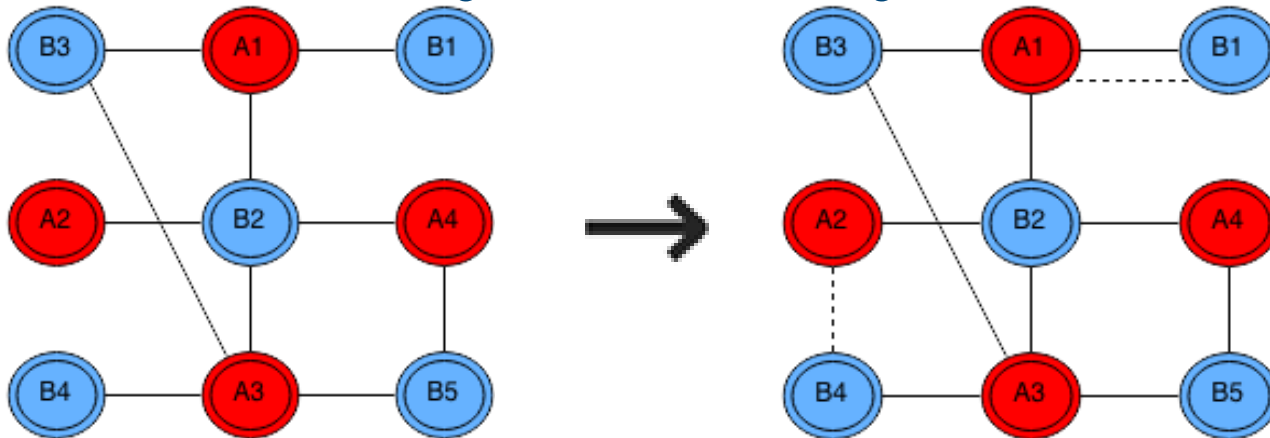


Pregel river, Königsberg

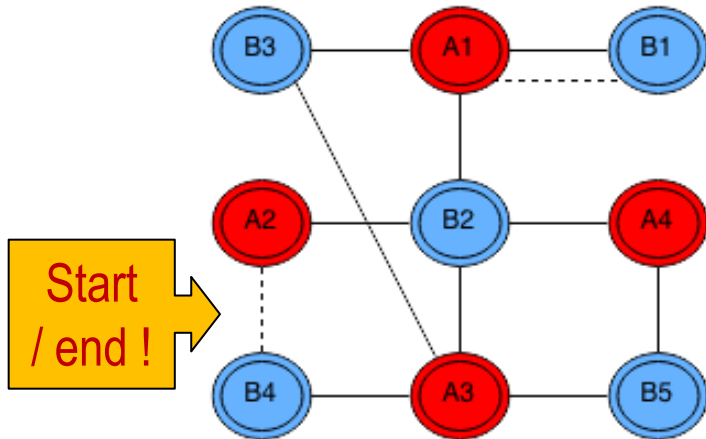


[Wikipedia]

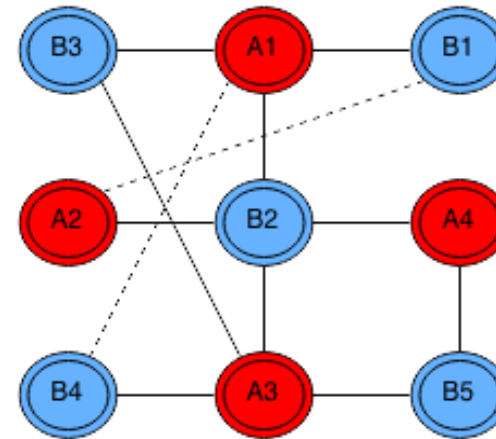
- Carl Hierholzer: „not even degree? Add aux edges!“



Finding Shortest Paths



<B4 , A3 , B2 , A1 , B1 , B3 ,
A3 , B5 , A4 , B2 , A2>

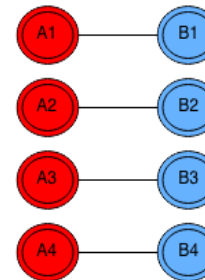
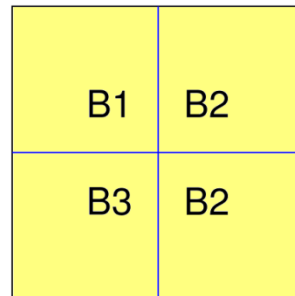
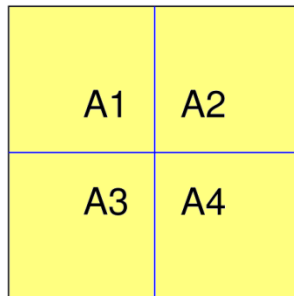


<B4 , A3 , B3 , A1 , B1 , A2 ,
B2 , A4 , B5 , A3 , B2 , A1>

- Assumption: can hold 1 red, 1 blue tile
- How to find **shortest** path? See paper!

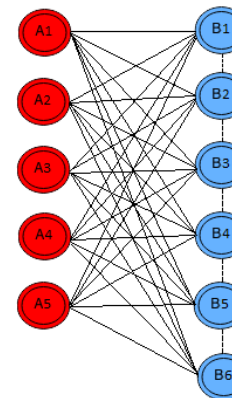
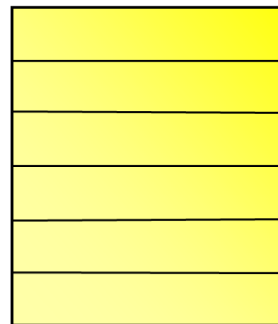
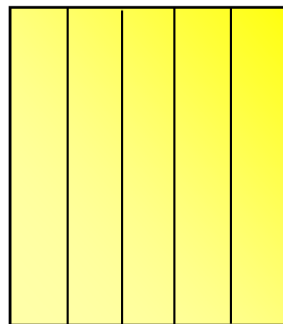
Complexity?

- Best case: full alignment



→ $|E_A| + |E_B|$

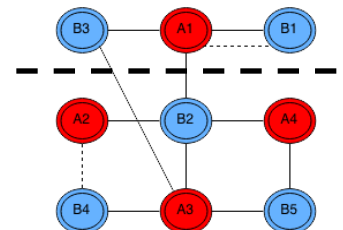
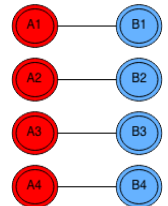
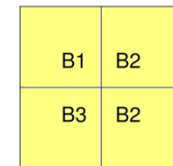
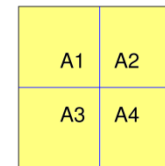
- Worst case: All tiles of A x all tiles of B



→ $|E_A| * |E_B|$

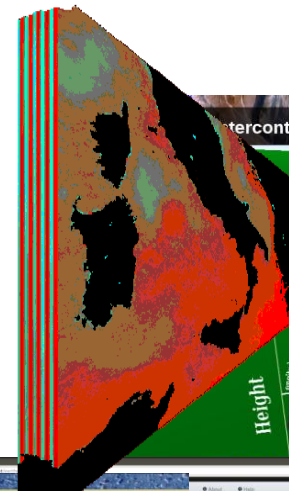
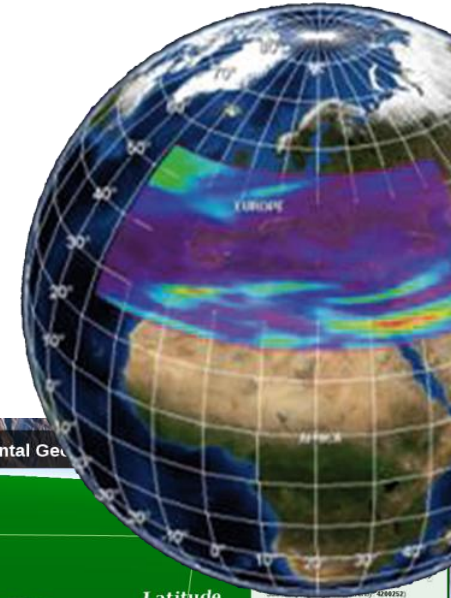
What else?

- We have: *tile traversal with minimum duplicate tile reads*
- Variation 1: *buffer size to avoid dup reads?*
 - query cost estimation
- Variation 2: *for buffer size given, how much improvement?*
- Variation 3: *parallelize disconnected subgraph*
- Variation 4: *how many tiles to ship between nodes for optimal parallelization?*

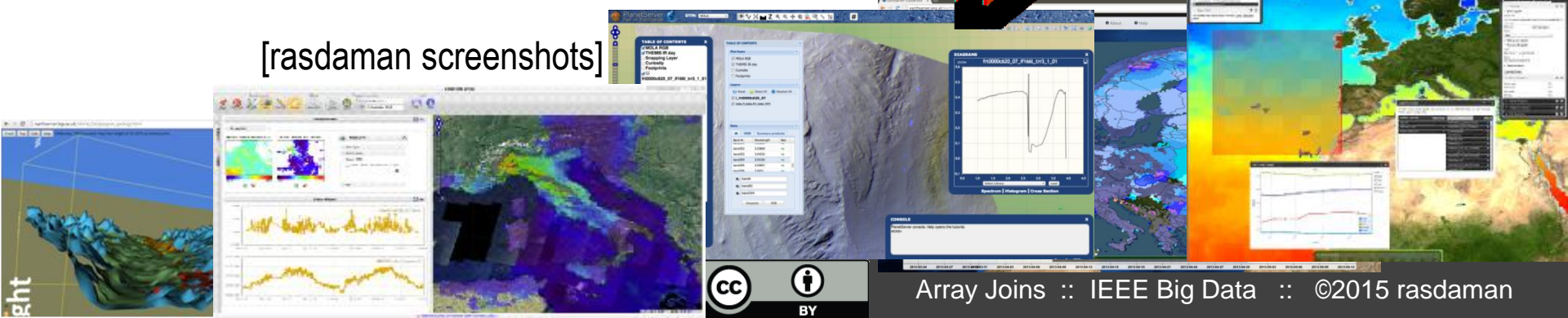


Wrap-Up

- Array Database queries offer new Big Data service level, including **datacube fusion**
 - Standardization: ISO Array SQL, OGC WCPS
- Need efficient algorithms
 - graph-based **Array Join** for arbitrary partitioning
- EarthServer: Petascale **Datacube Analytics**
 - rasdaman



[rasdaman screenshots]



Datacube Research @ Jacobs U

- Large-Scale Scientific Information Systems research group
 - Flexible, scalable n-D array services
 - www.jacobs-university.de/lis
- Main results:
 - pioneer Array DBMS, rasdaman
 - standardization:
 - OGC Big Geo Data (also ISO, INSPIRE, W3C)
 - ISO Array SQL

Hiring PhD students, PostDocs

